

# 10/10 - Percolation I

Tuesday, October 10, 2006  
10:51 AM

Are given a graph, typically infinite,  
a probability  $p$ , and a distinguished node, " $0$ "

Keep each edge with prob  $p$ .  
Call kept edges "open" and others "closed"

In infinite case, ask for prob that  $0$   
is in an infinite component

In finite case, ask for prob that  $0$   
is connected to boundary,  
or that there is a large component  
containing a constant fraction of the  
vertices.

Will begin with infinite case, as is easier.

Two infinite graphs:

$$\text{Lattice } \mathbb{Z}^2: \quad V = \mathbb{Z} \times \mathbb{Z}$$
$$E = \{(a,b), (c,d) : |a-c| + |b-d| = 1\}$$

distinguish  $(0,0)$

Infinite binary tree:  $T_2$

$$V = \{0,1\}^* \leftarrow \text{finite sequences of 0's and 1's}$$
$$E = (x,y) \in V \text{ s.t. } \text{length}(y) = \text{length}(x) + 1$$
$$\text{and } y = x0 \text{ or } y = x1$$

distinguish node  $\phi$  to empty set

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Let  $C$  be component of distinguished node.

$$\text{Define } \Theta(p) = \Pr_p[|C| = \infty]$$

the prob that  $C$  is infinite

Critical threshold is  $p_c$  st.

$$p_c = \sup \{ p \mid \Theta(p) = 0 \}$$

Is a phase change threshold

Can show that

$$p < p_c \quad \Pr[\exists \text{ an inf cluster}] = 0$$

$$p > p_c \quad \Pr[\exists \text{ an inf cluster}] = 1$$

By Kolmogorov's 0-1 Law

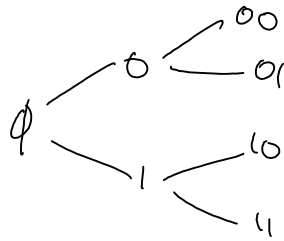
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For complete binary tree

$$\Pr[\phi \text{ is in inf component}]$$

$$= \Pr \left[ \left( \text{edge } (\phi, 0) \text{ appears } \wedge 0 \text{ in inf comp} \right) \vee \right. \\ \left. \left( \text{edge } (\phi, 1) \text{ appears } \wedge 1 \text{ in inf comp} \right) \right]$$

By self-similarity,  $\Pr[1 \text{ in inf comp}] = \Pr[\phi \text{ in inf comp}]$



So, by independence, and inclusion-exclusion

$$\Theta(p) = 2p\Theta(p) - (p\Theta(p))^2$$

So, either  $\Theta(p) = 0$ , or

$$1 = 2p - p^2\Theta(p) \Rightarrow \Theta(p) = \frac{2p-1}{p^2}$$

So, for  $p \leq \frac{1}{2}$ , is no solution

and, for  $p > \frac{1}{2}$ ,  $\Theta(p) > 0$ .  $\Rightarrow P_c = \frac{1}{2}$

Let's treat as limit of finite trees

Let  $T_2^d$  be binary tree of depth  $d$ .

Consider  $\Theta^d(p) = \text{prob is path from } \phi \text{ to a leaf in } T_2^d$

Consider  $\lim_{d \rightarrow \infty} \Theta^d(p)$  should =  $\Theta(p)$

By union bound,

$$\Theta^d(p) = \Pr[\text{is path from } \phi \text{ to a node in } \{0,1\}^d]$$

$$\begin{aligned}
&\leq \sum_{v \in \{0,1\}^d} \Pr[\text{is a path from } \phi \text{ to } v] \\
&= \sum_{v \in \{0,1\}^d} p^d = 2^d p^d = (2p)^d \rightarrow 0 \\
&\quad \text{if } p < \frac{1}{2}
\end{aligned}$$

So,  $P_c \leq \frac{1}{2}$

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To prove  $P_c \geq \frac{1}{2}$ ,

Note  $\theta^0(p) = 1$ , as just one node

$$\theta^1(p) = 2p - p^2$$

$$\theta^d(p) = 2p\theta^{d-1}(p) - (p\theta^{d-1}(p))^2$$

what is  $\lim_{d \rightarrow \infty} \theta^d(p)$ ?

Claim: if  $\theta^{d-1}(p) \geq \frac{2p-1}{p^2}$ , then

$$\theta^d(p) > \frac{2p-1}{p^2}$$

So, approaches limit  $> 0$

pf of claim

Consider  $f(x) = 2px - (px)^2$

We know  $f(x^*) = x^*$  for  $x^* = \frac{2p-1}{p^2}$

and  $f'(x) = 2p - 2p^2x > 0$  for  $0 \leq x \leq 1$   
so,  $f$  is monotone increasing.

$\Rightarrow$  for  $x > x^*$   $f(x) > f(x^*)$

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## Regular Grid

is known that  $P_c = \frac{1}{2}$ . Not easy

Will prove  $\frac{1}{3} \leq P_c \leq \frac{2}{3}$

Thm  $\frac{1}{3} \leq P_c$

Def  $\sigma(n) = \#$  loop-free walks from  $\bar{0}$   
of length  $n$ ,  $S(n) =$  set of loop-free paths  
of length  $n$

Prop  $\sigma(n) = 4 \cdot 3^{n-1}$

pf of Thm

If  $0$  is in inf component, then is a path  
from the origin of every length.

$$\begin{aligned} \theta(p) &\leq \Pr \left[ \text{there is a path } P \in S(n) \text{ all of whose edges open} \right] \\ &\leq \sum \Pr \left[ \text{all edges in } P \text{ open} \right] \end{aligned}$$

$$\begin{aligned}
 & \sum_{p \in S(n)} p^n \\
 &= \sum_{p \in S(n)} p^n = |S(n)| p^n \leq 4 \cdot 3^{n-1} p^n \\
 & \rightarrow 0 \text{ if } p < \frac{1}{3}
 \end{aligned}$$

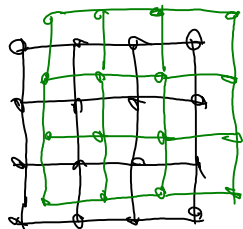
So  $\theta(p) = 0$  if  $p < \frac{1}{3}$ .

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Then  $p < \frac{2}{3}$ .

Will use dual graph:

vertex in center of every grid square,  
dual edges cross primal edges



Set dual edge open  $\Leftrightarrow$  corresp primal edge closed.

Prop

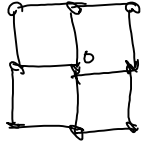
$0$  is not in inf cluster

$\Uparrow$   
Dual contains open cycle around  $0$ .

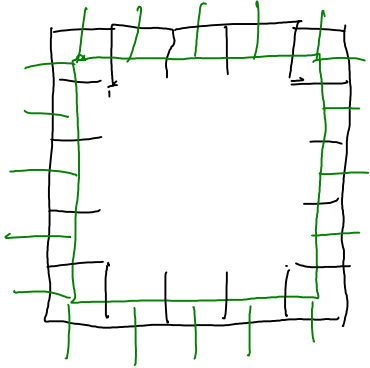
Let  $B_m = \text{box of edges on } [-m, m] \times [-m, m]$

$X_m = \text{event all edges in } B_m \text{ are open}$

$Y_m =$  event Dual does not contain open cycle around  $B_m$



$X_m$  all these present



$Y_m$  depends on these

$X_m \sim Y_m \Rightarrow \emptyset$  is a component

Let's show this probability is non-zero.

First, are independent, and  $\Pr[X_m] > 0$

So, just need to show  $\Pr[Y_m] > 0$

Will focus on  $\Pr[\bar{Y}_m] < 1$

Let  $a_n = \#$  cycles around origin of length  $n$  (simple).

How many? First, must pass through  $(\frac{1}{2}, \frac{1}{2} + k)$

for some  $0 \leq k \leq n$ .

From there, is a self-avoiding walk of

length  $n - 1$  (cut to least step)

$$\text{So, } |C(u)| \leq n \cdot S(n-1) \leq 4n \cdot 3^{n-2}$$

A cycle around  $B_m$  has length at least  $8m$

$$\text{So, } \Pr[\bar{Y}_m] \leq \sum_{n \geq 8m} \sum_{C \in C(u)} \Pr[\text{all dual edges of } C \text{ open}]$$

$$= \sum_{n \geq 8m} |C(u)| \cdot p^n \leq \sum_{n \geq 8m} 4n \cdot 3^{n-1} p^n$$

$$= \sum_{n \geq 8m} \left(\frac{4}{3}\right) n (3p)^n \rightarrow 0 \text{ for } n \text{ suff large}$$

$$\left( = \frac{4}{3} \left[ \frac{\binom{8m}{8m} (3p)^{8m}}{(1-3p)} + \frac{\binom{8m+1}{8m+1} (3p)^{8m+1}}{(1-3p)^2} \right] \right)$$