

Comparing Graphs. Lower bounds on λ_2 .

$A \succeq 0$ if A is positive semidefinite:
symmetric with no negative eigenvalues.

$$\Leftrightarrow x^T A x \geq 0, \forall x.$$

$$A \succeq B \text{ iff } A - B \succeq 0.$$

Is Loewner partial order.

$$A \succeq B \text{ and } B \succeq C \Rightarrow A \succeq C.$$

And for all symmetric C

$$A \succeq B \Rightarrow A + C \succeq B + C$$

Overload: define for graphs.

$$G \succeq H \text{ iff } L_G \succeq L_H$$

$$\text{Recall } L_G = \sum_{a \sim b} w_{a,b} (x(a) - x(b))^2$$

So, if H is a subgraph of G , $H \preceq G$
Same if decrease weights to get H .

Often consider $G \succeq cH$ where $c > 0$,

cH is H , but mult edge sets by c

$$L_{cH} = cL_H$$

Theorem If $G \geq cH$ then

$$\lambda_k(G) \geq c \lambda_k(H) \text{ for all } k$$

Proof by CF

$$\lambda_k(G) = \min_{\dim(S)=k} \max_{x \in S} \frac{x^T L_G x}{x^T x}$$

$$\geq \min_{\dim(S)=k} \max_{x \in S} c \frac{x^T L_H x}{x^T x}$$

$$= c \lambda_k(H)$$

Path inequality: $(n-1)P_n \geq G_{1,n}$

P_n is path from 1 to n . $G_{1,n}$ just has edge $(1,n)$

Proof $\forall x \in \mathbb{R}^n$, need to show

$$(n-1) \sum_{a=1}^{n-1} (x(a+1) - x(a))^2 \geq (x(n) - x(1))^2$$

$$\text{Set } \Delta(a) = x(a+1) - x(a), \text{ so } x(n) - x(1) = \sum_{a=1}^{n-1} \Delta(a).$$

$$\text{It's } (n-1) \sum_{a=1}^{n-1} \Delta(a)^2 \geq \left(\sum_{a=1}^{n-1} \Delta(a) \right)^2. \text{ It's Cauchy Schwarz.}$$

$$\left(\sum_{a=1}^{n-1} \Delta(a) \right)^2 = \left(\mathbf{1}_{n-1}^T \Delta \right)^2 = \|\mathbf{1}_{n-1}\|^2 \|\Delta\|^2 = (n-1) \sum_{a=1}^{n-1} \Delta(a)^2$$

Now, let's see how to use this to prove a lower bound on $\lambda_2(P_n)$.

last class saw $\lambda_2(P_n) \approx \frac{\pi^2}{n^2}$ and $\lambda_2(P_n) \leq \frac{12}{n(n+1)}$

To lower bound will prove

$P_n \geq cK_n$, and recall $\lambda_2(K_n) = n$,
so implies $\lambda_2(P_n) \geq cn$

Write $K_n = \sum_{a < b} G_{a,b}$

For $a < b$, write $P_{a,b}$ for subgraph of P_n from a to b

$$G_{a,b} \leq (b-a)P_{a,b} \leq (b-a)P_n$$

$$\Rightarrow K_n = \sum_{a < b} G_{a,b} \leq \sum_{a < b} (b-a)P_n = \frac{n(n+1)(n-1)}{6} P_n$$

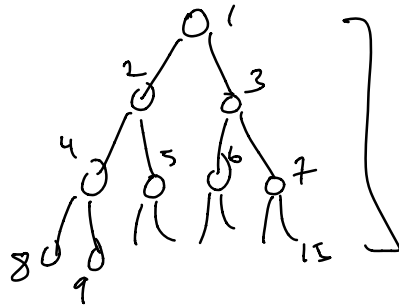
so, $P_n \geq \frac{6}{n(n+1)(n-1)}$

$$\Rightarrow \lambda_2(P_n) \geq n \cdot \frac{6}{(n+1)(n-1)n} = \frac{6}{(n+1)(n-1)}$$

Complete binary tree.

T_n

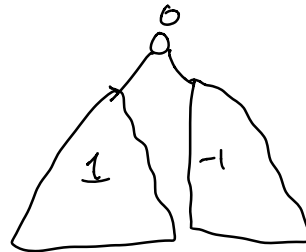
$$n = 2^{d+1} - 1$$



edges $a \rightarrow 2a$
 $a \rightarrow 2a+1$

depth 3

Upper bound on $\lambda_2(T_n)$ by



$$\frac{x^T L x}{x^T x} = \frac{2}{n-1} \geq \lambda_2$$

Lower bound:

For each $a < b$, let $T_n^{a,b}$ be the path in T_n from a to b .

Is unique because is a tree, and has length $\leq 2d$
 $\leq 2 \lg_2 n$

$$K_n = \sum_{a < b} G_{a,b} \leq \sum_{a < b} (2d) T_n^{a,b} \leq \sum_{a < b} (2d) T_n$$

$$\leq \binom{n}{2} (2 \lg_2 n) T_n = n(n-1) \lg_2 n T_n$$

$$\Rightarrow \lambda_2(T_n) \geq \frac{1}{n(n-1) \lg_2 n} \lambda_2(K_n) = \frac{1}{(n-1) \lg_2 n}$$

Differs from test vector bound by $\lg_2 n$

Sometimes use many paths.

Sometimes weight them.

Let P be the weighted path with edges $(a, a+1)$ of wt w_a .

$$\text{Then } G_{\text{un}} \preceq \left(\sum_{a=1}^{n-1} \frac{1}{w_a} \right) P$$

Proof For $x \in \mathbb{R}^n$, let $\Delta(a) = x(a+1) - x(a)$.

$$\text{Set } z(a) = \Delta(a) \sqrt{w_a}$$

$w^{-1/2}$ be vector s.t. $w^{-1/2}(a) = w(a)^{-1/2}$

$$\text{Then } \sum_a \Delta(a) = z^T w^{-1/2}$$

$$\begin{aligned} \text{so } x^T G_{\text{un}} x &= \left(\sum_a \Delta(a) \right)^2 = \left(z^T w^{-1/2} \right)^2 \\ &= \|w^{-1/2}\|^2 \|z\|^2 \\ &= \left(\sum_{a=1}^{n-1} \frac{1}{w(a)} \right) \cdot \sum_a w_a \Delta(a)^2 \\ &= \left(\sum_{a=1}^{n-1} \frac{1}{w(a)} \right) \cdot x^T L_P x \end{aligned}$$

Remind Is resistance of resistors in series, with $r(a) = \frac{1}{w(a)}$.

Approximations of graphs.

H is a c -approximation of G if

$$\frac{1}{c}H \preceq G \preceq cH$$

Next lecture: a random graph, every edge prob $\frac{1}{2}$ is a good $(1 \pm \epsilon)$, small ϵ -approximation of K_n

Expanders: For every $\epsilon > 0$, a constant d s.t. there is a d -regular $(1 \pm \epsilon)$ -approx of K_n

Ramanujan: $\epsilon \approx \frac{2}{\sqrt{d}}$

Sparsifiers: For every $\epsilon > 0$, is a constant d s.t. every G on n vertices has a $(1 \pm \epsilon)$ -approx with $\frac{dn}{2}$ (weighted) edges.